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Uncertainty and asset markets in GET

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- GET is capable of incorporating uncertainty albeit a little awkwardly.
- We focus on a pure exchange economy.
- The way to do this is to introduce states of the world $S = \{s_1, ..., s_R\}.$
- Consumer h evaluates that state s_i comes about with probability π_{hi}.
- The basic innovation is to regard, say, a loaf of bread in two different states as two different commodities.
- This means that consumption bundles and allocations are indexed by the states and become really long.

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Definition

A state contingent commodity vector/bundle $x = (x_{11}, ..., x_{L1}, ..., x_{1R}, ..., x_{LR}) \in \mathbb{R}^{LR}_+$ provides a consumer bundle $(x_{1s}, ..., x_{Ls}) \in \mathbb{R}^{L}_+$ in state $s \in S$.

- Notice that above is a description of a contingent bundle for a consumer.
- An allocation, a bundle for each consumer, is an element of \mathbb{R}_+^{LRH}
- Consumer *h* has endowment $\omega_h = (\omega_{11h}, ..., \omega_{L1h}, ..., \omega_{1Rh}, ..., \omega_{LRh}).$
- Preferences are defined over contingent commodity=bundles. = -> ...

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• The consumer has a Bernoulli-utility function for each state and evaluates bundles by weighing the states by their probabilities so that $x_h \succeq_h x'_h$ if and only if

$$\sum_{s \in S} \pi_{sh} u_{sh}(x_{1sh}, ..., x_{Lsh}) \geq \sum_{s \in S} \pi_{sh} u_{sh}(x'_{1sh}, ..., x'_{Lsh})$$

- Formally, the economy with states and state contingent commodities is equivalent to the standard version, and all the results we have there remain true.
- In particular, Walras-equilibria exist and they are Pareto-efficient.
- When the Bernoulli-utility functions are concave we immediately see that in equilibrium there is efficient risk sharing.

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Information

- Think of information becoming more and more accurate with time, and to model this assume that there are dates $t \in \{0, 1, ..., T\}$.
- A partition of a set A is a collection of sets $\{A_i\}_{i=1}^n$ such that $A_i \cap A_j = \emptyset$ when $i \neq j$, and $\bigcup_{i=1}^n A_i = A$.
- When A = S the subsets of S, $S_i = A_i$, are called events, and a partion \mathscr{L} of S is called an information structure.
- When s, s' ∈ S_i then an agent regards it as possible that the state of the world is s as well as s'.
- Information revelation is modelled by a sequence of information structures (L₀, L₁,..., L_T) such that if Q ∈ L_t then there exists R ∈ L_{t-1} such that Q ⊆ R. (B) < E) = 0.0

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• Each consumer could have his/her own information structure but to keep notation from exploding we assume that they have a common sequence of information structures.

• A pair (t, E) where $E \in \mathscr{L}_t$ is called a date-event.

• There are *L* commodities and now we denote contingent commodities by index *It* meaning that commodity *I* is available at date *t*.

 We still need to index the commodities by states so x_{lts} is the amount of commodity *l* available at date *t* in state *s* (along with other states that belong to the same event).

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 We can define the double-indexed commodities, of which there are L(T+1), as our new commodities, and then everything is as in the standard setting.

 The only thing we have to take care is that bundles and allocations are measurable, i.e., if s, s' ∈ S_i then x_{lts} = x_{lts'}.

• This way the above temporal structure can be made atemporal, and the economy is formally like the standard setting.

• The equilibrium of this economy is called the Arrow-Debreu equilibrium.

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Definition

Arrow-Debreu equilibrium consists of allocation $\begin{aligned} x^* &= (x_1^*, x_2^*, ..., x_H^*) \in \mathbb{R}_+^{LRH}, \\ x_h^* &= (x_{11h}^*, ..., x_{L1h}^*, ..., x_{1Rh}^*, ..., x_{LRh}^*) \in \mathbb{R}_+^{LR}, \text{ and a system of prices} \\ p &= (p_{11}, ..., p_{L1}, ..., p_{1R}, ..., p_{LR}) \in \mathbb{R}_+^{LR} \text{ such that} \\ \text{i) for every } h \ x_h^* \text{ is the maximal element in the budget set} \\ \left\{ x_h \in \mathbb{R}_+^{LR} : px_h \leq p\omega_h \right\} \text{ and} \\ \text{ii) markets clear } \sum_{h=1}^H x_h^* = \sum_{h=1}^H \omega_h. \end{aligned}$

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Radner-equilibrium

- In the Arrow-Debreu economy there are markets for all contingent commodities, and trading happens before any of the uncertainty is resolved.
- Assume that at t = 0 the agents can trade just one commodity but then at the succeeding dates there are spotmarkets at each state.
- Now only S forward markets exist (one for each state) in the one commodity that is traded at t = 0.
- To proceed we need to postulate that the agents have (correct) expectations about the spot prices at time t = 1.
- The expected price at state s is denoted by $p_s = (p_{1s}, ..., p_{Ls})$ and the expected vector of prices is $p = (p_{s_1}, ..., p_{s_R})$.
- At time t = 0 there is contingent trade in good 1, and its prices are given by $q = (q_{s_1}, ..., q_{s_R})$.

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- Given the prices consumer h makes a consumption plan for contingent commodities (z_{s1h},..., z_{sRh}) ∈ ℝ^R and for spot markets (x_{s1h},..., x_{sRh}) where one has to be careful with notation as x_{sh} = (x_{1sh},..., x_{Lsh}) ∈ ℝ^L₊.
- The consumer's problem is given by

$$\max_{\substack{\left(z_{s_1h},...,z_{s_Rh}\right) \in \mathbb{R}^R \\ \left(x_{s_1h},...,x_{s_Rh}\right) \in \mathbb{R}_+^{LR}} U_h(x_{s_1h},...,x_{s_Rh})$$

$$s.t. \quad \frac{\sum_{s \in S} q_s z_{sh} \le 0}{p_s x_{sh} \le p_s \omega_{sh} + p_{1s} z_{sh}} \quad \forall s \in S$$

- The first budget constraint does not restrict much; it is well possible that, say, $z_{s_1h} < -\omega_{1s_1h}$.
- In this case the consumer sells short and has to acquire sufficient amount of good 1 in the spot market to honour his/her commitments.
- The requirement of positive consumption (i.e., wealth) at each state guarantees that things do not get out of control.

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Definition

A Radner-equilibrium consists of prices at time t = 0 for contingent first good commodities $q = (q_{s_1}, ..., q_{s_R})$ and spot prices $p = (p_{s_1}, ..., p_{s_R})$ at time t = 1, for each consumer h consumption plan $(z_{s_1h}^*, ..., z_{s_Rh}^*)$ at time t = 0, and consumption plan $(x_{s_1h}^*, ..., x_{s_Rh}^*)$ at time t = 1 such that the consumption plans solve the above maximisation problem, and $\sum_{h \in \mathscr{H}} z_{sh}^* \leq 0$ and $\sum_{h \in \mathscr{H}} x_{sh}^* \leq \sum_{h \in \mathscr{H}} \omega_{sh}$ for each $s \in S$.

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Theorem

i) Assume that an Arrow-Debreu equilibrium is given by allocation $x^* \in \mathbb{R}^{LRH}_+$ and a contingent commodity prices $p = (p_{s_1}, ..., p_{s_R}) \in \mathbb{R}_{++}^{LR}$. There are prices $q = (q_{s_1}, ..., q_{s_R}) \in \mathbb{R}_{++}^{R}$ for the contingent commodity-1 and consumption plans $z^* = (z_1^*, ..., z_H^*) \in \mathbb{R}^{RH}$ such that the consumption plans x^* , z^* and prices g and p constitute a Radner-equilibrium. ii) If prices $q = (q_{s_1}, ..., q_{s_P}) \in \mathbb{R}_{++}^R$ for the contingent commodity-1 and spot prices $p = (p_{s_1}, ..., p_{s_R}) \in \mathbb{R}_{++}^{LR}$, and consumption plans $z^* = (z_1^*,...,z_H^*) \in \mathbb{R}^{RH}$ and $x^* \in \mathbb{R}^{LRH}_+$ constitute a Radner-equilibrium then there are multipliers $(\mu_{s_1}, ..., \mu_{s_R}) \in \mathbb{R}^R_{++}$ such that allocation x^* and the contingent commodities prices $(\mu_{s_1}p_{s_1},...,\mu_{s_R}p_{s_R}) \in \mathbb{R}_{++}^R$ constitute an Arrow-Debreu equilibrium.

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Proof.

i) A good guess is to make $q_s = p_{1s}$ for each $s \in S$. Let us next compare the budget sets of consumer h under Arrow-Debreu prices and Radner-prices. The Arrow-Debreu budget set is $B_h^{AD} = \left\{ (x_{s_1h}, ..., x_{s_Rh}) \in \mathbb{R}_+^{LR} : \sum_{s \in S} p_s (x_{sh} - \omega_{sh}) \leq 0 \right\}$ while the Radner budget set is

$$B_h^R =$$

$$\left\{ (x_{s_1h}, \dots, x_{s_Rh}) \in \mathbb{R}_+^{LR} : \exists (z_{s_1h}, \dots, z_{s_Rh}) \in \mathbb{R}^R \ s.t.$$
$$\sum_{s \in S} q_s z_{sh} \le 0 \ and \ p_s (x_{sh} - \omega_{sh}) \le p_{1s} z_{sh}, \ \forall s \in S \right\}$$

Assume that $x_h \in B_h^{AD}$, and let $z_{sh} = \frac{1}{p_{1s}} p_s (x_{sh} - \omega_{sh})$. Now $\sum_{s \in S} q_s z_{sh} = \sum_{s \in S} p_s (x_{sh} - \omega_{sh}) \le 0$ and $p_s (x_{sh} - \omega_{sh}) = p_{1s} z_{sh}$ for all $s \in S$. This shows that $B_h^{AD} \subseteq B_h^R$.

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Proof.

Assume that $x_h \in B_h^R$. This means that there exists $(z_{s_1h}, ..., z_{s_Rh})$ such that $\sum_{s \in S} q_s z_{sh} \leq 0$ and $p_s (x_{sh} - \omega_{sh}) \leq p_{1s} z_{sh}$ for all $s \in S$. Summing the latter inequality over the states yields $\sum_{s \in S} p_s (x_{sh} - \omega_{sh}) \leq \sum_{s \in S} p_{1s} z_{sh} = \sum_{s \in S} q_s z_{sh} \leq 0$. This shows that $B_h^R \subseteq B_h^{AD}$.

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Proof.

ii) Choose $\mu_s=rac{q_s}{p_{1s}}$ for all $s\in S$, and write the Radner budget set as $B_{h}^{R} =$ $\left\{ (x_{s_1h}, \dots, x_{s_Rh}) \in \mathbb{R}^{LR}_+ : \exists (z_{s_1h}, \dots, z_{s_Rh}) \in \mathbb{R}^R \, s.t. \right\}$ $\sum_{s \in S} q_s z_{sh} \leq 0 \text{ and } \mu_s p_s (x_{sh} - \omega_{sh}) \leq p_{1s} z_{sh}, \forall s \in S$ $s \in S$ But then one can mimic the proof of i) and write the budget set in

the Arrow-Debreu form. QED

- Assets pay money or real goods conditional on states.
- Let us assume that asset payments are in good 1.
- There are two dates t = 0 and t = 1.
- Let there be S states like in 19.E in MWG.
- Asset is characterised by a return vector $r = (r_1, ... r_S) \in \mathbb{R}^S$.
- r_s is what the asset pays in state s.

- Options constitute an important class of assets.
- Underlying an option is another asset; assume its return vector is given by $r \in \mathbb{R}^{S}$.
- A call option is characterised by a strike price c.
- A unit of an option gives the holder a right to buy a unit of the underlying asset at price *c* after the uncertainty has resolved.
- The return vector of the asset is $r(c) = (max \{0, r_1 - c\}, max \{0, r_2 - c\}, ..., max \{0, r_5 - c\}).$
- At state s an option holder exercises his/her right if $r_s c > 0$.

- Assume that there is a given set of assets called asset structure tradable at t = 0.
- Assume there are K assets.
- The price of assets at t=0 is given by $q=(q_1,...,q_K)\in\mathbb{R}^K$.
- The trades are denoted by $z = (z_1, ..., z_K) \in \mathbb{R}^K$, and called a portfolio.

Asset markets

Definition

Asset prices $q = (q_1, ..., q_K) \in \mathbb{R}^K$ and spot prices $p_s = (p_{1s}, ..., p_{Ls}) \in \mathbb{R}^L$ for all s, and for all consumers i a portfolio $z_i^* = (z_{1i}^*, ..., z_{Ki}^*) \in \mathbb{R}^K$, and consumption plan $x_i^* = (x_{1i}^*, ..., x_{Si}^*) \in \mathbb{R}^{LS}$ is a Radner-equilibrium if z_i^* and x_i^* solve

$$max_{z_i^* \in \mathbb{R}^K x_i^* \in \mathbb{R}^{LS}} U_i(x_{1i}, ..., x_{Si})$$

$$s.t.\sum_k q_k z_k \leq 0$$

$$p_s x_{si} \leq p_s \omega_{si} + \sum_k p_{1s} z_{ki} r_{sk} \forall s$$

$$\sum_{i}^{2} x_{ki}^{*} \leq 0$$
$$\sum_{i} x_{si}^{*} \leq \sum_{i} \omega_{si} \forall s, k$$

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- Note that $p_{1s} = 1$ is a possible normalisation; this is done in the sequel.
- Things are clearer in matrix-notation and we define the $S \times K$ return matrix

where row s keeps track of returns of different assets in state s, and column k keeps track of asset k's returns in different states.

• Now the budget constraint becomes

$$B_i(p,q,R) = \begin{cases} p_1(x_{1i} - \omega_{1i}) & \vdots \\ x \in \mathbb{R}_+^{LS} : z_i \in \mathbb{R}^K, qz_i \leq 0 \text{ and } & \vdots \\ p_5(x_{5i} - \omega_{5i}) & \vdots \end{cases} \leq Rz_i$$

Theorem

Assume that $0 \neq r_k \geq 0$ for all k. Then for every vector of asset prices $q \in \mathbb{R}^K$ in a Radner-equilibrium can be found multipliers $\mu = (\mu_1, ..., \mu_S) \geq 0$ such that $q_k = \sum_s \mu_s r_{sk}$, or $q^T = \mu R$.

Proof.

Asset price vector $q \in \mathbb{R}^{K}$ is arbitrage free if there is no portfolio $z = (z_1, ..., z_K)$ such that $qz \leq 0, 0 \neq Rz \geq 0$. It is straightforward that in equilibrium q must be arbitrage free. We first show that if $q \in \mathbb{R}^{K}$ is arbitrage free then there exists a vector of multipliers $\mu = (\mu_1, ..., \mu_S) \geq 0$ such that $q^T = \mu R$. It is clear that $q_k > 0$ for all k. We can also assume that each row of R has strictly positive elements.

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Asset markets

Proof.

follows

Consider the set

$$V = \left\{ v \in \mathbb{R}^{S} : v = Rz, z \in \mathbb{R}^{K} \text{ and } qz = 0 \right\}$$

Since q is arbitrage free $V \cap \mathbb{R}^S_+ \setminus \{0\} = \emptyset$. Both are convex and there exists a separating hyperplane; there exists $\mu' = (\mu'_1, ..., \mu'_S) \ge 0$ such that $\mu' v \le 0$ for $v \in V$ and $\mu' w \ge 0$ for $w \in \mathbb{R}^S_+$. Since $-v \in V$ whenever $v \in V$ it must be the case that

$$\mu' v = 0$$
 for all $v \in V$

We claim that q^T is proportional to $\mu' R \in \mathbb{R}^K$. First $0 \neq \mu' R \geq 0^T$. If the claim does not hold there exists $\overline{z} \in \mathbb{R}^K$ such that $q\overline{z} = 0$ and $\mu' R\overline{z} > 0$. But if $v = R\overline{z}$ then $v \in V$ and $\mu' v \neq 0$ which is a contradiction. Consequently, $q^T = \alpha \mu' R$ for some $\alpha > 0$; let $\mu = \alpha \mu'$. As equilibrium asset prices must be arbitrage free the result

Definition

Asset structure with $S \times K$ return matrix is complete if rank R = S.

- With complete asset structure transfer of wealth over all states is possible.
- If S = 4 and a primary asset has returns r = (4,3,2,1) we can have a complete asset structure with options r(3.5), r(2.5) and r(1.5).

Theorem

If the asset structure is complete the equilibrium is Pareto-efficient.

• With assets the important thing is the range of the return matrix

range
$$R = \left\{ v \in \mathbb{R}^{S} : v = Rz, z \in \mathbb{R}^{K} \right\}$$

 It tells the possible wealth vectors that are achievable with a given asset structure.

Theorem

Let there be two assets structures with return matrices R and R'. If range R = range R' then the consumption plans x^* in Radner equilibria corresponding to R and R' are equal.

Exercise 19.E.4 Suppose that $r_3 = \alpha_1 r_1 + \alpha_2 r_2$. Show that in equilibrium $q_3 = \alpha_1 q_1 + \alpha_2 q_2$. Assume first that $q_3 > \alpha_1 q_1 + \alpha_2 q_2$. Then portfolio $z = (\alpha_1 q_3, \alpha_2 q_3, -\alpha_1 q_1 - \alpha_2 q_2)$ must return zero or qz = 0. Now $\sum_{k} p_{1s} r_{sk} z_{k} = p_{1s} (r_{s1} \alpha_{1} q_{3} + r_{s2} \alpha_{2} q_{3} - r_{s3} (\alpha_{1} q_{1} + \alpha_{2} q_{2})) =$ $p_{1s}r_{s3}(q_3-(\alpha_1q_1+\alpha_2q_2))$ for each s. Since $0 \neq r_3 > 0$ we have $\sum_{k} p_{1s} r_{sk} z_k \geq 0$ for each s and at least one strict inequality. Consumers can always increase their wealth, and utility, by adding zto their portfolio. But this cannot happen in equilibrium. If $q_3 < \alpha_1 q_1 + \alpha_2 q_2$ then we can show analogously that consumers can always profitably subtract z from their portfolio. Thus, it must be the case that $q_3 = \alpha_1 q_1 + \alpha_2 q_2$.

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